


Let I be a small category, i.e. the objects of I form a set

Def: A diagram of shape I in a category \mathcal{C} is a functor $F: I \rightarrow \mathcal{C}$

Ex: $I =$ 

$F(1) = X_1$
 $f \downarrow \downarrow g$
 $F(2) = X_2$

$\text{Fun}(I, \mathcal{C}) =$ functor category, i.e

obj: functors $I \xrightarrow{F} \mathcal{C}$

mor. natural transformations:

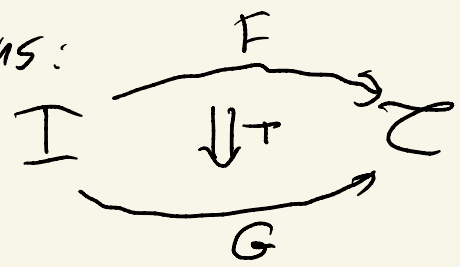
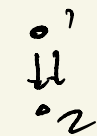
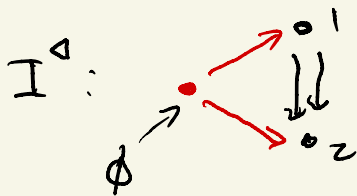


diagram category.

Def: left cone of \mathcal{I} : $\mathcal{I}^\triangleleft \leftarrow$ category

$$\text{ob}(\mathcal{I}^\triangleleft) = \text{ob}(\mathcal{I}) \sqcup \{\phi\}$$

Ex: $\mathcal{I} =$ 



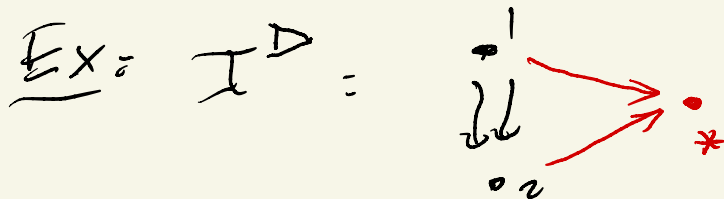
↑
is an initial object in $\mathcal{I}^\triangleleft$
all other morphisms are the morphisms in \mathcal{I}

Def: right cone : $\mathcal{I}^\triangleright$

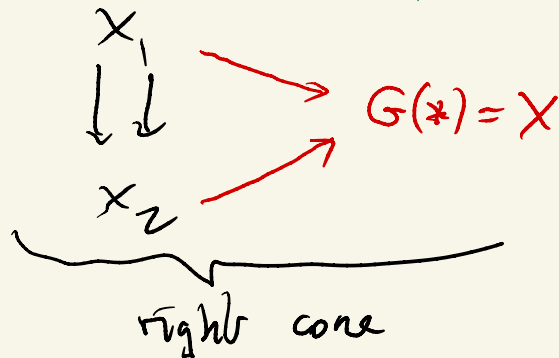
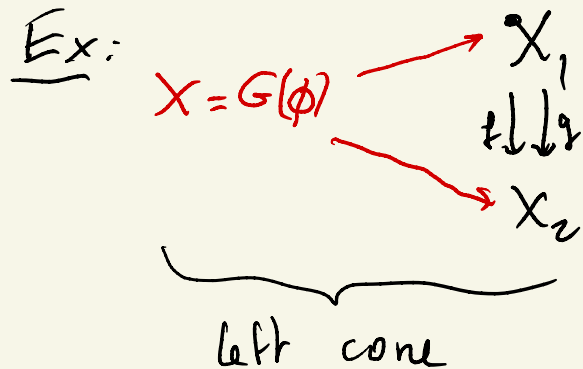
$$\text{ob}(\mathcal{I}^\triangleright) = \text{ob}(\mathcal{I}) \sqcup \{*\}$$

↑ terminal object in $\mathcal{I}^\triangleright$

$\mathcal{I} \subset \mathcal{I}^\triangleright$
full subcat

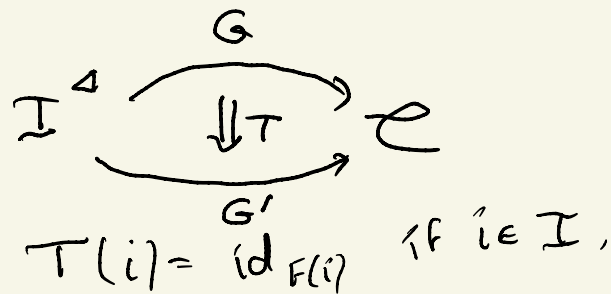


Def: A left cone ^(right) over a diagram $F: I \rightarrow \mathcal{C}$
 is a functor $G: I^\Delta \rightarrow \mathcal{C}$ s.t. $G|_I = F$
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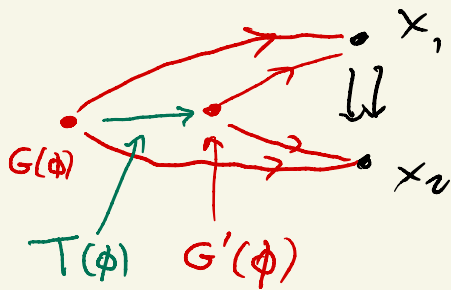


Def: left cones over F are the objects in the category $\text{lcone}(F)$.

morphisms $G \rightarrow G'$:



Ex:

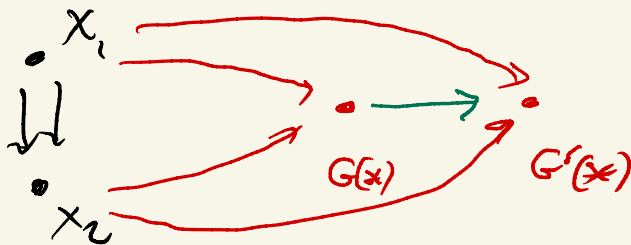


$$\begin{array}{ccc}
 y \xrightarrow{T} x & \begin{array}{c} \gamma, x \in \mathcal{I}^{\Delta} \\ G(y) \rightarrow G(x) \in \mathcal{C} \end{array} & \begin{array}{c} G(\phi) \\ \downarrow \\ G'(y) \rightarrow G'(x) \in \mathcal{C} \\ G'(\phi) \end{array} \\
 & & \begin{array}{c} \downarrow T(y) \quad \square \quad \downarrow T(x) \end{array}
 \end{array}$$

naturality of T means that all these diagrams commute.

similarly, we define the category $\text{rcone}(F)$ of right cones over F .

Ex:



Def. $F: I \rightarrow \mathcal{C}$ is a diagram in \mathcal{C}

A limit diagram is a terminal object G in the category $\mathcal{L}(\text{Cone}(F))$.

Then $G(\phi)$ is called the limit of F ,

$\lim_{I \triangleleft} F$ notation: $G(\phi) = \lim_I F$

or $\lim(F: I \rightarrow \mathcal{C})$

or $\lim_{i \in I} F(i)$



$\lim_I F$ is well defined up to unique isomorphisms.

A colimit diagram is an initial object G in $\text{rCone}(F)$.

$G(x) \in \mathcal{C}$ is called the colimit of F over I .

Notation: $\text{colim}_I F$ or $\text{colim}_{i \in I} F(i)$ or $\text{colim}(F: I \rightarrow \mathcal{C})$.

Def: If all limits in \mathcal{C} exist, then \mathcal{C} is called complete.

if all limits in \mathcal{C} exist, then \mathcal{C} is called cocomplete.

Explicit construction in $\mathcal{C} = \text{Set, Top}$:

$X: I \rightarrow \text{Set}$

$$\lim_I X = \left\{ (x_i) \in \prod_{i \in I} X(i) \mid \begin{array}{l} f: i \rightarrow j, \text{ i.e. } f \in I(i, j) \\ \left. \begin{array}{l} X(i) \xrightarrow{X(f)} X(j) \\ \underbrace{x_i}_{x_i} \longmapsto X(f)(x_i) = x_j \end{array} \right\} \end{array} \right\}$$

↑ requirement
of morphisms f

$$\text{colim}_{\underline{I}} X = \frac{\coprod_{i \in \underline{I}} X(i)}{\sim} \quad \begin{array}{l} \swarrow \text{equiv. relation} \\ \text{gen. by} \end{array} \quad \begin{array}{l} \phi \in \underline{I}(i, j) \\ X(i) \xrightarrow{X(\phi)} X(j) \end{array}$$

in particular Set, Top are complete + cocomplete.

for $\mathcal{C} = \text{Vect}$: limit description is as for Set, Top .

$$V: \underline{I} \rightarrow \text{Vect} \quad \text{colim}_{\underline{I}} V = \frac{\bigoplus_{i \in \underline{I}} V(i)}{\sim} \quad \begin{array}{l} \swarrow \text{equiv. relation} \\ \text{gen. by} \end{array} \quad \begin{array}{l} \phi \in \underline{I}(i, j) \\ V(i) \xrightarrow{V(\phi)} V(j) \end{array}$$

Rem: if \underline{I} has an initial object $\phi \in \underline{I}$ then $\lim_{\underline{I}} F \cong F(\phi)$ HW!
 $\ast \in \underline{I}$ terminal object $\Rightarrow \text{colim}_{\underline{I}} F \cong F(\ast)$

$V: \mathcal{I} \rightarrow \text{Vect}$	$\lim_{\mathcal{I}} V$	$\text{colim}_{\mathcal{I}} V$
\mathcal{I} discrete category, i.e. only identity morphisms.	$\prod_{i \in \mathcal{I}} V(i)$	$\bigoplus_{i \in \mathcal{I}} V(i)$
$V_1 \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{g} \end{array} V_2$	$\text{eq}(V_1 \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{g} \end{array} V_2)$	$\text{coeq}(V_1 \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{g} \end{array} V_2)$
$U \begin{array}{c} \xrightarrow{f} \\ \downarrow g \\ W \end{array} V$	U	$V \oplus W / \begin{array}{l} f(u) \sim g(u) \\ \text{for all } u \in U \end{array}$ <u>pushout</u>
$W \begin{array}{c} \xrightarrow{g} \\ \downarrow f \\ U \end{array} V$	$\{(v, w) \in V \times W \mid f(v) = g(w)\}$ <u>pullback</u>	U
$\rho: G \rightarrow \text{Aut}(W)$ representation $\mathcal{X}/G = \mathcal{I} \rightarrow \text{Vect}$ <u>obj</u> : one object $\mathcal{X} \mapsto W$ <u>mor</u> $(\mathcal{X}, \mathcal{X}) \cong G \ni g \mapsto \rho(g)$	$W^G = \{w \in W \mid \rho(g)w = w\}$ G -invariants	$W_G = W / \{w \sim \rho(g)w\}$ G -coinvariants.