

Advances in Interval Methods for Deterministic Global Optimization in Chemical Engineering

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4th International Conference on Frontiers in Global Optimization
Santorini, Greece
June 8–12, 2003

Outline

- Background
- Interval Methodology
- Solving the Linear Interval Subproblem
- Examples
 - Parameter estimation problems
 - Problems with many local minima
 - Stationary states for triatomic molecules
- Concluding Remarks

Background

- One approach for doing global optimization is the use of **interval analysis**.
- Interval analysis can:
 - Provide a **deterministic** methodology for global optimization problems.
 - Deal automatically with rounding error, thus providing both **mathematical and computational guarantees**.
- Interval methods can be used in various ways in global optimization, e.g.:
 - Interval branch-and-bound
 - ⇒ **Interval-Newton approach**
 - As a tool within other methods

Background (Cont'd)

- Interval Newton/Generalized Bisection (IN/GB)
 - Given a system of equations to solve, an initial interval (bounds on all variables), and a solution tolerance:
 - IN/GB can find (enclose) with **mathematical and computational certainty** either **all** solutions or determine that no solutions exist.
 - IN/GB can also be extended and employed as a deterministic approach for global optimization problems.
 - A general-purpose approach; in general requires no simplifying assumptions or problem reformulations.
 - No strong assumptions about functions need to be made.
- ⇒ **Solution of a linear interval equation system is a key subproblem.**

Interval Methodology (Cont'd)

Problem: Solve $\mathbf{f}(\mathbf{x}) = \mathbf{0}$ for all roots in interval $\mathbf{X}^{(0)}$.

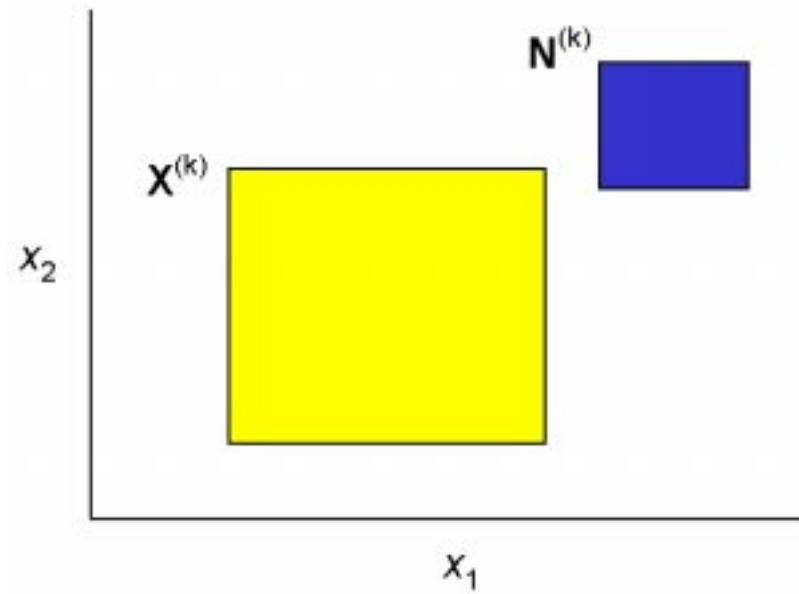
Basic iteration scheme (IN/GB): For a particular subinterval (box), $\mathbf{X}^{(k)}$, perform root inclusion test:

- (Range Test) Compute an interval extension (bounds on range) for each function in the system.
 - If $\mathbf{0}$ is not an element of any interval extension, delete the box. Otherwise,
- (Interval-Newton Test) Compute the *image*, $\mathbf{N}^{(k)}$, of the box by solving the linear interval equation system

$$\mathbf{F}'(\mathbf{X}^{(k)})(\mathbf{N}^{(k)} - \tilde{\mathbf{x}}^{(k)}) = -\mathbf{f}(\tilde{\mathbf{x}}^{(k)})$$

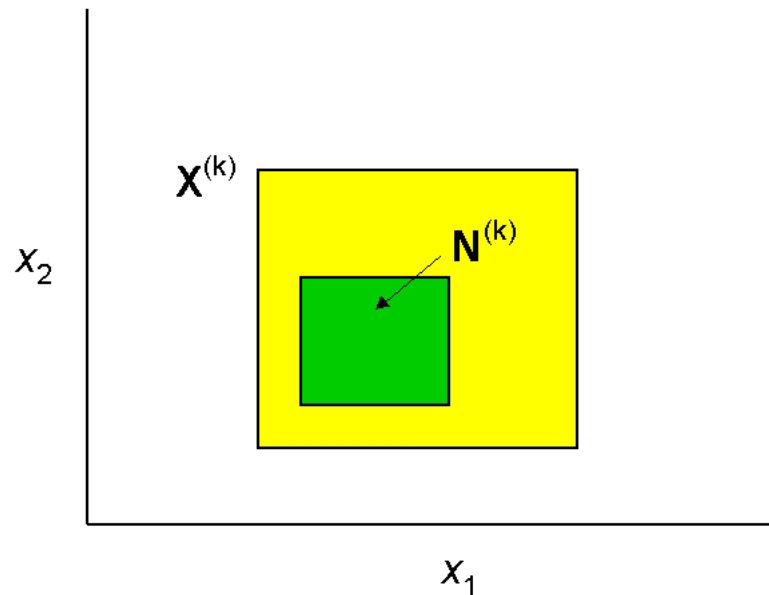
- $\tilde{\mathbf{x}}^{(k)}$ is some point in $\mathbf{X}^{(k)}$.
- $\mathbf{F}'(\mathbf{X}^{(k)})$ is an interval extension of the Jacobian of $\mathbf{f}(\mathbf{x})$ over the box $\mathbf{X}^{(k)}$.

Interval Methodology (Cont'd)



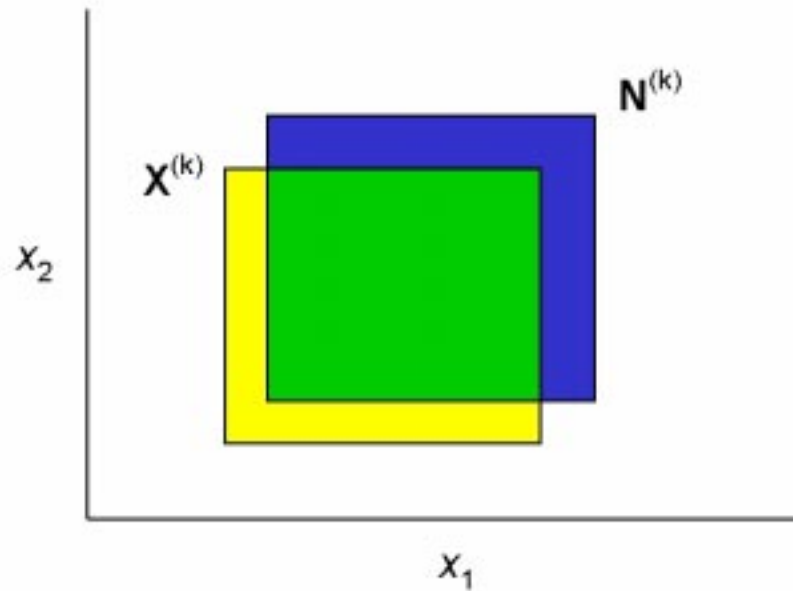
- There is no solution in $X^{(k)}$.

Interval Methodology (Cont'd)



- There is a *unique* solution in $\mathbf{X}^{(k)}$.
- This solution is in $\mathbf{N}^{(k)}$.
- Additional interval-Newton steps will tightly enclose the solution with quadratic convergence. (Point Newton method will also converge to solution from any point in $\mathbf{N}^{(k)}$.)

Interval Methodology (Cont'd)



- Any solutions in $\mathbf{X}^{(k)}$ are in intersection of $\mathbf{X}^{(k)}$ and $\mathbf{N}^{(k)}$.
- If intersection is sufficiently small, repeat root inclusion test.
- Otherwise, bisect the intersection and apply root inclusion test to each resulting subinterval.

Interval Methodology (Cont'd)

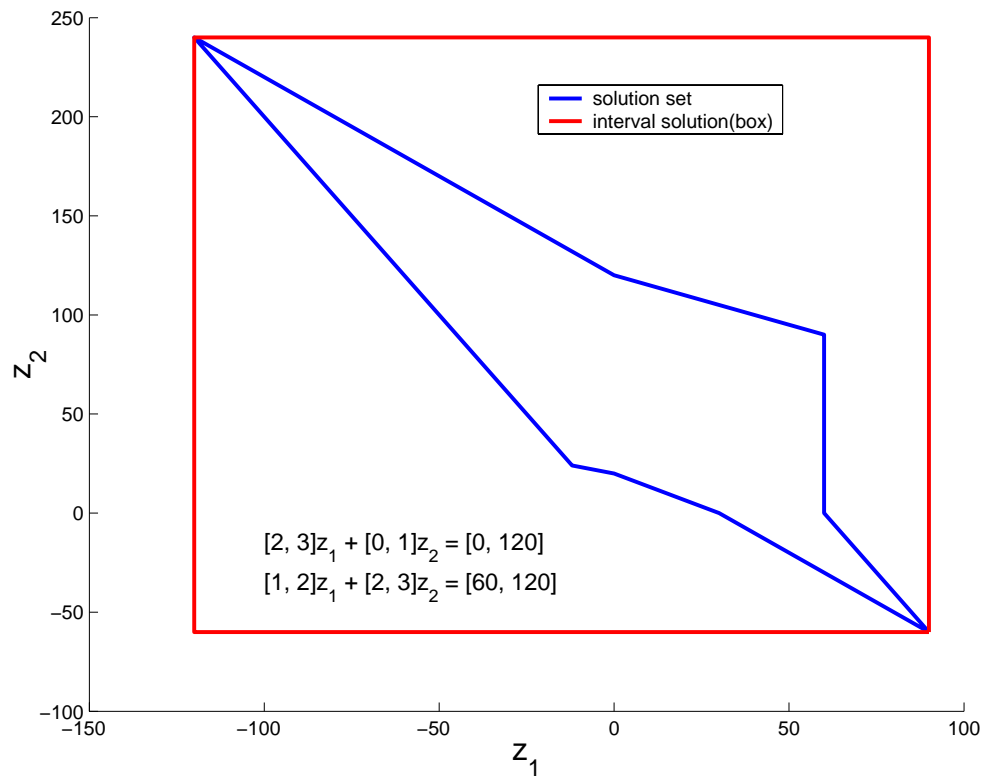
- Easily extended to **global optimization** problems.
- For unconstrained problems, solve for **stationary points**.
- For constrained problems, solve for **KKT** or **Fritz-John** points.
- Add an additional pruning condition (objective range test):
 - Compute interval extension of objective function.
 - If its lower bound is greater than a known upper bound on the global minimum, prune this subinterval.
- This combines IN/GB with a branch-and-bound scheme.
- Key step, for either optimization or equation solving, is solution of linear interval system

$$\mathbf{F}'(\mathbf{X})(\mathbf{N} - \tilde{\mathbf{x}}) = -\mathbf{f}(\tilde{\mathbf{x}})$$

Seek tightest possible bounds on solution $(\mathbf{N} - \tilde{\mathbf{x}})$, and thus on \mathbf{N} .

Solution Set of Linear Interval System

- Consider linear interval system $\mathbf{Az} = \mathbf{B}$.
- Solution set is defined: $\mathbf{S} = \{\mathbf{z} \mid \tilde{\mathbf{A}}\mathbf{z} = \mathbf{b}, \tilde{\mathbf{A}} \in \mathbf{A}, \mathbf{b} \in \mathbf{B}\}$.
- Interval solution: An interval \mathbf{Z} containing \mathbf{S} .



Solution Set of Linear Interval System (Cont'd)

- Computing the interval hull (tightest interval containing \mathbf{S}) is NP-hard (Rohn and Kreinovich, 1995).
 - Several methods are available to compute an interval solution \mathbf{Z} that contains \mathbf{S} , but that may not give tight bounds.
 - Methods used in the context of interval-Newton:
 - Preconditioned (inverse-midpoint) interval Gauss-Seidel
 - Hybrid (pivoting/inverse-midpoint) preconditioner and real point selection (HP/RP) (Gau and Stadtherr, 2002)
- ⇒ **LP strategy**

LP Strategy for Linear Interval System

- Oettli & Prager(1964) theorem : Solution set \mathbf{S} is defined by the constraints

$$\left| \hat{\mathbf{A}}\mathbf{z} - \hat{\mathbf{B}} \right| \leq \Delta\mathbf{A} |\mathbf{z}| + \Delta\mathbf{B}$$

$\hat{\mathbf{A}}$ – component-wise midpoint matrix of \mathbf{A}

$\Delta\mathbf{A}$ – component-wise half width matrix of \mathbf{A}

$\hat{\mathbf{B}}$ – component-wise midpoint vector of \mathbf{B}

$\Delta\mathbf{B}$ – component-wise half width vector of \mathbf{B}

- To eliminate absolute value operation on \mathbf{z} , the components of \mathbf{z} must keep a constant sign \longrightarrow **consider each orthant separately.**

LP Strategy for Linear Interval System (Cont'd)

- In each orthant, define D_α , a diagonal matrix whose entries are:

$$(D_\alpha)_{jj} = \begin{cases} 1 & \mathbf{z}_j \geq 0 \\ -1 & \mathbf{z}_j < 0 \end{cases} \quad j = 1, 2, \dots, n$$

- To determine bounds on \mathbf{S} in each orthant, solve $2n$ linear programming problems:

$$\begin{aligned} & \text{maximize (and minimize) } \mathbf{z}_j, \quad j = 1, 2, \dots, n \\ \text{s.t. } & \begin{pmatrix} \hat{\mathbf{A}} - \Delta\mathbf{A}D_\alpha \\ -\hat{\mathbf{A}} - \Delta\mathbf{A}D_\alpha \end{pmatrix} \mathbf{z} \leq \begin{pmatrix} \bar{\mathbf{B}} \\ -\underline{\mathbf{B}} \end{pmatrix} \end{aligned}$$

- To get optimal solution overall (interval hull), calculate extrema in all orthants (2^n in worst scenario — **exponential complexity**).

LP Strategy for Linear Interval System (Cont'd)

Application to IN/GB methods:

- Solve linear interval system

$$\mathbf{F}'(\mathbf{X})(\mathbf{N} - \tilde{\mathbf{x}}) = -\mathbf{f}(\tilde{\mathbf{x}})$$

- Only the part of \mathbf{N} that intersects \mathbf{X} needs to be found.
- If $\tilde{\mathbf{x}}$ is selected to be a corner of \mathbf{X} , then the part of $\mathbf{N} - \tilde{\mathbf{x}}$ for which \mathbf{N} lies in \mathbf{X} is entirely in one orthant.
- Solution of interval-Newton equation can be sought using LP in only one orthant. **Tightest possible** solution obtained, while **avoiding exponential time complexity**.

Numerical Experiments

- LISS_LP(Linear Interval System Solver by Linear Programming) has been developed.
- Option to use sparse linear algebra in solution of LP problem.
- We compare performance results of LISS_LP to HP/RP (Gau and Stadtherr, 2002) on a [SUN Blade 1000 model 1600 workstation](#).
- Performance results include:
 - Number of interval Newton tests performed (I-N tests)
 - CPU time in seconds

Example 1

- Estimation of Van Laar parameters from vapor-liquid equilibrium data using error-in-variables approach (Kim et al., 1990; Esposito and Floudas, 1998):

$$P = \gamma_1 x_1 p_1^0(T) + \gamma_2 (1 - x_1) p_2^0(T)$$

$$y_1 = \frac{\gamma_1 x_1 p_1^0(T)}{\gamma_1 x_1 p_1^0(T) + \gamma_2 (1 - x_1) p_2^0(T)}$$

where

$$p_1^0(T) = \exp \left[18.5875 - \frac{3626.55}{T - 34.29} \right], \quad p_2^0(T) = \exp \left[16.1764 - \frac{2927.17}{T - 50.22} \right]$$

and

$$\gamma_1 = \exp \left[\frac{A}{RT} \left(1 + \frac{A}{B} \frac{x_1}{1 - x_1} \right)^{-2} \right], \quad \gamma_2 = \exp \left[\frac{B}{RT} \left(1 + \frac{B}{A} \frac{1 - x_1}{x_1} \right)^{-2} \right]$$

- There are five data points and four **measured variables** with two **parameters** to be determined.

Example 1 (Cont'd)

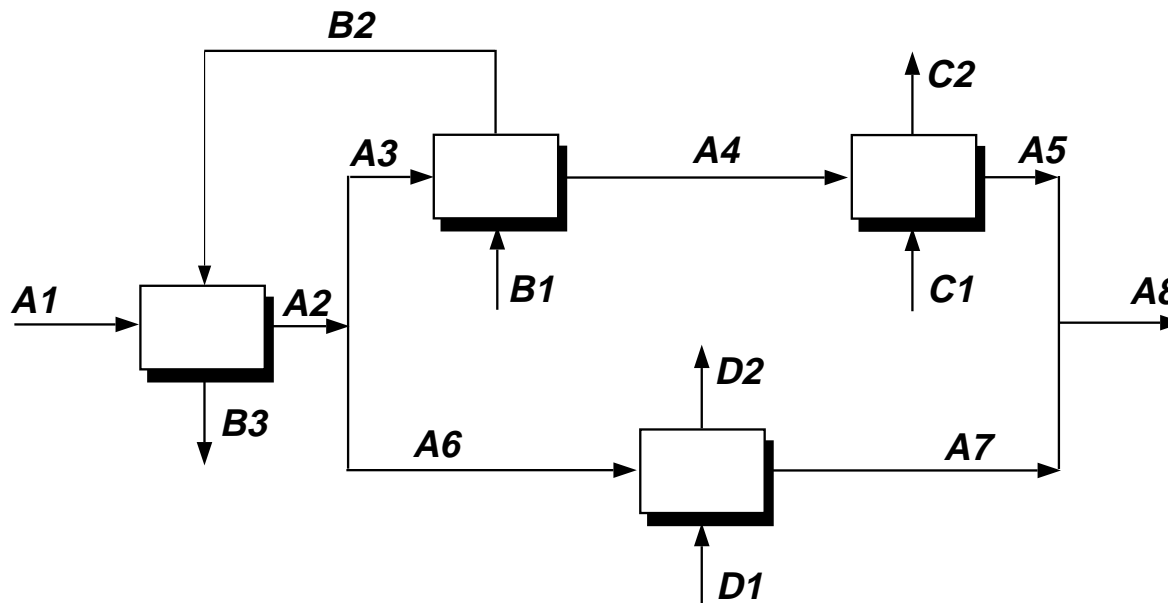
- Formulated as unconstrained global optimization with 2 parameter variables and 10 state variables.
- With standard inverse-midpoint preconditioner in solution of linear interval system, solution time is > 2 CPU days.
- Performance results with hybrid preconditioner (HP/RP) and LISS_LP:

	HP/RP	LISS_LP
I-N tests	303,589	156,182
CPU time (s)	664.4	496.7

- LP solver uses dense linear algebra.

Example 2

- Estimation of parameters in heat exchanger network using error-in-variables approach (Biegler and Tjoa, 1993).
- Network of four exchangers. Estimate the four rating parameters $(UA)_i$.
- Five unconstrained global optimization problems with 4 parameter variables and $13m$ state variables (number of data points $m = 4, 8, 12, 16, 20$).



Example 2 (Cont'd)

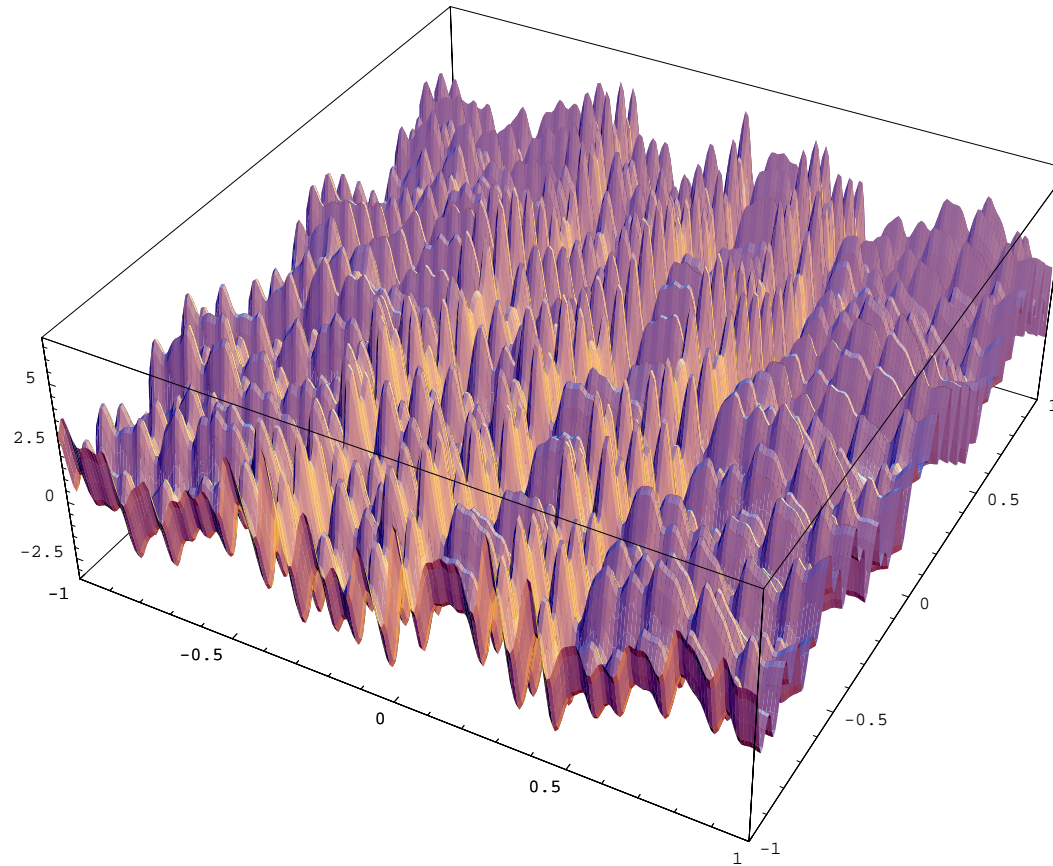
- Performance results:

Data Points	Variables	HP/RP		LISS_LP	
m	n	I-N Tests	CPU Time	I-N Tests	CPU Time
4	56	1	0.12	2	0.27
8	108	375	211.8	44	38.1
12	160	363	498.6	299	346.0
16	212	188	645.8	83	316.8
20	264	220	1357.3	81	504.9

- LP solver uses sparse linear algebra.

Example 3

- Trefethen (2002) Challenge Problem #4 — Find the Global Minimum



$$f(x, y) = \exp(\sin(50x)) + \sin(60 \exp(y)) + \sin(70 \sin(x)) + \sin(\sin(80y)) - \sin(10(x + y)) + (x^2 + y^2)/4; \quad x \in [-1, 1]; \quad y \in [-1, 1]$$

Example 3 (Cont'd)

- Solution

$$x \in [-0.02440307969437517, -0.02440307969437516]$$

$$y \in [0.2106124271553557, 0.2106124271553558]$$

$$f \in [-3.306868647475245, -3.306868647475232]$$

- Global minimum is easily found using interval approach

	HP	LISS_LP
I-N tests	1814	1179
CPU time (s)	0.15	0.16

- On relatively easy problems, LP-based strategy is not needed, but still can be used without significant loss of efficiency due to LP overhead.

Example 4

- Find the global minimum of the function (Siirola et al., 2002):

$$f(\mathbf{x}) = 100 \prod_{i=1}^N \sum_{j=1}^5 \left(\frac{j^5}{4425} \cos(j + jx_i) \right) + \frac{1}{N} \sum_{i=1}^N (x_i - x_{0,i})^2$$

where $x_{0,i} = 3$, $x_i \in [x_{0,i} - 20, x_{0,i} + 20]$, $i = 1, \dots, N$.

- Solve for $N = 5$. There are $\approx 10^8$ local optima.
- Performance results:

	Global Minima	I-N tests	CPU time(s)
LISS_LP	5	155,666	389.17
HP	5	171,918	636.76

Example 5

- Find all stationary points (minima, maxima, saddles) on potential energy surface of triatomic molecule ABC.
 - Useful for study of transition states and reaction pathways.
 - We studied the molecules HCN, HSiN, CS₂ and HBO (Westerberg and Floudas, 1999).
- The triatomic molecule geometry is described using the three interatomic distances $R_1 = R_{AB}$, $R_2 = R_{AC}$ and $R_3 = R_{BC}$.
- Murrel-Sorbie analytic potential energy surface $V(R_1, R_2, R_3)$ is used.
- Find **all** solutions of stationarity condition: $\nabla V = 0$

Example 5 (Cont'd)

- Physical constraints: triangle inequality

$$R_{AB} \leq R_{AC} + R_{BC}$$

$$R_{AC} \leq R_{AB} + R_{BC}$$

$$R_{BC} \leq R_{AB} + R_{AC}$$

Use to tighten interval bounds before application of interval-Newton.

- Collinear case: eliminate one variable, e.g., for $R_{AC} = R_{AB} + R_{BC}$ solve

$$\frac{\partial}{\partial R_{AB}} V(R_{AB}, R_{AB} + R_{BC}, R_{BC}) = 0,$$

$$\frac{\partial}{\partial R_{BC}} V(R_{AB}, R_{AB} + R_{BC}, R_{BC}) = 0$$

- Solve **one** noncollinear problem and **three** collinear problems.
- Search in intervals $R_i \in [0.7, 5.0] \text{Å}$.

Example 5 (Cont'd)

- General form of potential energy surface (Aguilar et al., 1992)

$$V(R_1, R_2, R_3) = V_{AB} + V_{AC} + V_{BC} + V_{ABC}$$

- Two-body terms are extended Rydberg functions, e.g.,

$$V_{AB}(R_1) = -D_e(1 + a_1\rho + a_2\rho^2 + a_3\rho^3 + \dots)e^{-a_1\rho},$$

$$\rho = R_1 - R_e$$

- Three-body term

$$V_{ABC}(R_1, R_2, R_3) = P \times T,$$

$$P = V^0 \left(1 + \sum_{i=1}^3 C_i \rho_i + \sum_{j \geq i=1}^3 C_{ij} \rho_i \rho_j + \sum_{k \geq j \geq i=1}^3 C_{ijk} \rho_i \rho_j \rho_k + \dots \right)$$

$$T = \prod_{i=1}^3 \left(1 - \tanh \frac{\gamma_i S_i}{2} \right)$$

$$S_i = \sum_{j=1}^3 b_{ij} \rho_j, \quad \rho_j = R_j - R_j^0$$

Example 5 (Cont'd)

- Stationary states for HCN

Type	Energy (eV)	$R_{CN}(\text{\AA})$	$R_{CH}(\text{\AA})$	$R_{NH}(\text{\AA})$
minimum	-5.548223	—	2.332871	1.038900
saddle	8.094668	—	0.857572	0.806900
minimum	-12.972507	1.159150	—	0.993336
minimum	-13.592215	1.153198	1.065498	—
saddle	-5.249952	2.344235	2.980408	1.044278
saddle	-1.937592	2.311895	1.792854	2.327696
saddle	-3.102483	2.582864	1.081559	2.737335
saddle	-11.444169	1.117973	1.053919	1.387750
saddle	-11.345398	0.929065	1.039138	1.041348
minimum	-11.379410	0.857321	0.980845	0.989052

Example 5 (Cont'd)

- Stationary states for HSiN

Type	Energy (eV)	$R_{SiN}(\text{Å})$	$R_{SiH}(\text{Å})$	$R_{NH}(\text{Å})$
saddle	1.109601	2.778074	2.617596	—
saddle	-3.144738	1.523964	2.426268	—
saddle	-5.148745	2.006322	1.361586	—
minimum	-6.098598	1.529588	1.459586	—
saddle	-5.666608	1.575921	—	2.969229
minimum	-9.358509	1.523293	—	0.998205
maximum	1.720515	2.647092	2.415995	3.498876
saddle	-2.876954	1.501907	2.309780	3.069649
saddle	-4.908995	2.394221	2.137984	0.974496
saddle	-0.728138	2.155741	1.473092	2.044809
saddle	-3.717494	1.461352	1.634575	2.093708

Example 5 (Cont'd)

- Stationary states for CS₂

Type	Energy (eV)	$R_{CS}(\text{\AA})$	$R_{CS'}(\text{\AA})$	$R_{SS'}(\text{\AA})$
saddle	-1.668827	2.761779	—	2.695109
saddle	103.740892	0.949956	—	1.813411
minimum	97.485407	0.909824	—	1.417728
minimum	-12.004548	1.552422	1.552422	—
saddle	-0.049002	4.171034	4.171034	3.978688

Example 5 (Cont'd)

- Stationary states for HBO(PES1)

Type	Energy (eV)	$R_{BH}(\text{\AA})$	$R_{BO}(\text{\AA})$	$R_{OH}(\text{\AA})$
saddle	-7.598281	3.264082	1.187662	—
minimum	-16.678316	1.165505	1.185028	—
minimum	-6.556670	1.162756	—	2.349430
saddle	-0.216647	—	2.554092	3.688901

- Stationary states for HBO(PES2)

Type	Energy (eV)	$R_{BH}(\text{\AA})$	$R_{BO}(\text{\AA})$	$R_{OH}(\text{\AA})$
minimum	-16.678851	1.168947	1.184167	—
minimum	-6.639249	1.154136	—	2.344208
minimum	-11.305022	—	1.192047	2.383483
saddle	-11.134196	2.906575	1.185979	2.398381

Example 5 (Cont'd)

- Summary of triatomic problems

Problem	Stationary points found	CPU time (sec)
HCN	10	6.66
HSiN	11	1.07
CS ₂	5	2.18
HBO(PES1)	4	0.88
HBO(PES2)	4	0.56

Concluding Remarks

- An LP-based method can be used to solve the linear interval system arising in the context of the interval-Newton approach for nonlinear equation solving and global optimization.
- The method can obtain tighter bounds on the solution set than standard methods, and thus lead to a large reduction in the number of subintervals that must be tested during the interval-Newton procedure.
- The overhead required to solve the LP subproblems may lead to relatively smaller improvements in overall computation time.
- The interval methodology is a powerful approach for deterministic global optimization.

Acknowledgments

- ACS Petroleum Research Fund 35979-AC9
- Indiana 21st Century Research & Technology Fund
- NSF DMR 00-79647