

# Membership tests for images of algebraic sets by linear projections

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## Abstract

Given a witness set for an irreducible variety  $V$  and a linear map  $\pi$ , we describe membership tests for both the constructible algebraic set  $\pi(V)$  and the algebraic set  $\overline{\pi(V)}$ . We also provide applications and examples of these new tests including computing the codimension one components of  $\overline{\pi(V)} \setminus \pi(V)$ . Additionally, we also describe computing the geometric genus of a curve section of an irreducible component of the solution set of a polynomial system and a test for deciding whether a plane quartic curve is a Lüroth quartic.

**Keywords.** Numerical algebraic geometry, polynomial system, algebraic sets, witness sets, projections, membership test, numerical irreducible decomposition, geometric genus, Lüroth quartic

**AMS Subject Classification.** 65H10, 68W30

## Introduction

Given a polynomial system  $f : \mathbb{C}^N \rightarrow \mathbb{C}^n$ , an  $\ell$ -dimensional irreducible component  $V \subset f^{-1}(0)$ , and a linear map  $\pi : \mathbb{C}^N \rightarrow \mathbb{C}^K$ , a “witness set” for  $\overline{\pi(V)}$  was constructed in [7] from a witness set for  $V$ , hereafter called a *pseudo-witness set* for  $\overline{\pi(V)}$ . This approach reduces computations on  $\overline{\pi(V)}$  to computations on  $V$  without using elimination theory to construct a polynomial system  $g$  such that  $\overline{\pi(V)}$  is an irreducible component of  $g^{-1}(0)$ .

The main results of this article, presented in §2, are algorithms for performing a numerical membership test for both  $\pi(V)$  and  $\overline{\pi(V)}$ .

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Chevalley's Theorem [4] states that the image of a constructible set, e.g.,  $\pi(V)$ , is a constructible set<sup>1</sup>. Effective symbolic methods for performing computations with constructible sets are discussed in [5, 19].

In §3, we use these membership tests to compute a numerical decomposition of the irreducible components of  $\overline{\pi(V)} \setminus \pi(V)$  of codimension one in  $\overline{\pi(V)}$  and use this to develop an approach for computing the geometric genus of a generic curve section of  $\overline{\pi(V)}$ .

The necessary background material is presented in §1 which also codifies the properties of our substitute for witness sets into the notion of a pseudo-witness set.

In §4, we present examples using our new membership tests.

## 1 Background material

We collect some background material in this section. Throughout, we assume  $f : \mathbb{C}^N \rightarrow \mathbb{C}^n$  is a polynomial system and define  $\mathcal{V}(f)$  to be the set of points in  $\mathbb{C}^N$  which  $f$  maps to 0. The algebraic set  $\mathcal{V}(f)$  is reduced and, in particular, all of the irreducible components of  $\mathcal{V}(f)$  have multiplicity one. We let  $f^{-1}(0)$  denote  $\mathcal{V}(f)$  with its underlying scheme structure, which includes the multiplicity information of the components of  $\mathcal{V}(f)$  with regard to  $f$ .

### 1.1 Witness sets

Suppose that  $V \subset f^{-1}(0)$  is an  $\ell$ -dimensional irreducible algebraic set of degree  $d$ . A *witness set* for  $V$  is the triple  $\{f, \mathcal{L}, W\}$  where  $\mathcal{L}$  consists of  $\ell$  general linear polynomials on  $\mathbb{C}^N$  and  $W = V \cap \mathcal{V}(\mathcal{L})$ . The *witness point set*  $W$  consists of  $d$  points. A finite set  $\mathcal{W}$  with  $W \subset \mathcal{W} \subset V$  is called a *witness point superset* for  $V$ .

The *multiplicity* of  $V$  with respect to  $f$  is the multiplicity of any  $w \in W$  as a root of  $\begin{bmatrix} f \\ \mathcal{L} \end{bmatrix}$ . The component  $V$  is said to be *generically reduced* with respect to  $f$  if the multiplicity of  $V$  with respect to  $f$  is 1. Otherwise,  $V$  is said to be *generically nonreduced*, which we consider in the following section. See [23, Chap. 13] for more details regarding witness sets.

### 1.2 Deflation

If  $V$  is generically nonreduced with respect to  $f$ , then the deflation approach of [10] produces a polynomial system  $F : \mathbb{C}^N \rightarrow \mathbb{C}^m$ , with  $m \geq n$ , such that  $F^{-1}(0)$  has an irreducible and generically reduced component  $\hat{V}$  which, as a set, is equal to  $V$ . By renaming as necessary, we will assume *without loss of generality* that  $V$  is generically reduced with respect to  $f$ .

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<sup>1</sup>A *constructible subset* of an algebraic set  $X$  is any set in the Boolean algebra of subsets of  $X$  obtained by starting with algebraic subsets of  $X$  and closing up under the operations of finite unions and complementation.

It should be noted that more traditional versions of deflation (see also [8, 11, 12] and [23, §13.3.2, §15.2.2]) change the dimension of the ambient space and may replace  $V$  with an algebraic set  $V'$  that maps generically one-to-one onto a dense subset of  $V$ .

### 1.3 Randomization

Let  $f : \mathbb{C}^N \rightarrow \mathbb{C}^n$  be a polynomial system and  $1 \leq k \leq n$ . For  $A \in \mathbb{C}^{k \times (n-k)}$ , let

$$\mathcal{R}(f; k) = [I_k \ A] \cdot f$$

where  $I_k$  is the  $k \times k$  identity matrix. It is a consequence of Bertini's theorem, e.g., [22] or [23, §13.5], that any irreducible codimension  $k$  component of  $\mathcal{V}(f)$  is an irreducible component of  $\mathcal{V}(\mathcal{R}(f; k))$  for a nonempty Zariski open (and hence dense) set of matrices  $A \in \mathbb{C}^{k \times (n-k)}$ . Thus, we will assume *without loss of generality* that  $f : \mathbb{C}^N \rightarrow \mathbb{C}^k$  is a polynomial system where  $V \subset f^{-1}(0)$  is a codimension  $k$  irreducible component.

### 1.4 Pseudo-witness sets

Let  $f : \mathbb{C}^N \rightarrow \mathbb{C}^n$  be a polynomial system and  $\{f, \mathcal{L}, W\}$  be a witness set for an irreducible and generically reduced component  $V \subset f^{-1}(0)$  of dimension  $\ell$ . Suppose that  $\pi : \mathbb{C}^N \rightarrow \mathbb{C}^K$  is a linear map and  $B \in \mathbb{C}^{K \times N}$  such that  $\pi(x) = Bx$ .

Even though the set  $\pi(V)$  might not be an algebraic set, it is very close to an algebraic set. More specifically,  $\pi(V)$  is a *constructible algebraic set* which means that it is a member of the Boolean algebra of sets constructed from algebraic sets by the operations of finite unions, finite intersections, and complementation. A typical example is the projection onto  $(x, y)$  of  $\mathcal{V}(x - yz)$ : the image is  $(\mathbb{C}^2 \setminus \mathcal{V}(y)) \cup \{(0, 0)\}$ .

The closure of a constructible algebraic set  $C$  in the complex topology  $\overline{C}$  is the same as the closure of  $C$  in the Zariski topology. The same statement holds for the interior  $C^\circ$  of  $C$  with  $\overline{C^\circ} = \overline{C}$ . In particular, since the dimensions of  $\overline{C}$  and  $C^\circ$  are equal, the dimension of  $C$  is well-defined. Finally, if  $\overline{C}$  is pure  $k$ -dimensional, then  $\overline{C} \cap L = C^\circ \cap L$  for a general affine linear space  $L$  of codimension  $k$ . Additional details for constructible algebraic sets is provided in [23, Appendix A].

Let  $\ell' = \dim \pi(V)$ . For  $i = 1, \dots, \ell'$ , let  $b_i \in \mathbb{C}^N$  be general elements in the row span of  $B$  and, for  $i = \ell' + 1, \dots, \ell$ , let  $b_i \in \mathbb{C}^N$  be general elements in  $\mathbb{C}^N$ . We call the quadruple  $\{f, \pi, \mathcal{L}', W'\}$  [7], where

$$\mathcal{L}'(x) = \begin{bmatrix} b_1 \cdot x - 1 \\ \vdots \\ b_{\ell'} \cdot x - 1 \end{bmatrix} \quad \text{and} \quad W' = V \cap \mathcal{V}(\mathcal{L}'),$$

a *pseudo-witness set* for  $\pi(V)$  with  $\deg \overline{\pi(V)} = |\pi(W)|$ .

A pseudo-witness set may be efficiently used to fulfill the same tasks for which a witness set for  $\overline{\pi(V)}$  would be used if we had a set of polynomials on  $\mathbb{C}^K$  whose solution set contained  $\overline{\pi(V)}$  as an irreducible component. One example is using pseudo-witness

sets in place of witness sets to work with the numerical irreducible decomposition [20] of the closure of the image of an algebraic map, e.g., [1, §2.1.3].

## 1.5 Moving linear spaces

The membership tests developed in this article are based on moving linear spaces. Let  $\{f, \mathcal{L}, W\}$  be a witness set for an irreducible and generically reduced  $V \subset f^{-1}(0)$  of dimension  $\ell$  where  $f : \mathbb{C}^N \rightarrow \mathbb{C}^{N-\ell}$  and  $\mathcal{L} : \mathbb{C}^N \rightarrow \mathbb{C}^\ell$ . Given a system of linear polynomials  $\widehat{\mathcal{L}} : \mathbb{C}^N \rightarrow \mathbb{C}^\ell$  with  $\dim \mathcal{V}(\widehat{\mathcal{L}}) = N - \ell$ , we want to compute the set of points  $\widehat{W} := V \cap \mathcal{V}(\widehat{\mathcal{L}}) \subset V$  by deforming  $\mathcal{L}$  to  $\widehat{\mathcal{L}}$  using the “square” homotopy  $H : \mathbb{C}^N \times \mathbb{C} \rightarrow \mathbb{C}^N$  defined by

$$H(x, t) = \begin{bmatrix} f(x) \\ (1-t)\widehat{\mathcal{L}}(x) + t\mathcal{L}(x) \end{bmatrix}. \quad (1)$$

Starting at  $t = 1$  with the points in  $W$ , continuation allows one to track the path defined by  $H(x, t) \equiv 0$  as  $t$  goes from 1 to 0. Additional details are provided in [23].

Of the  $|W|$  paths tracked using the homotopy  $H$ , some of them may diverge as  $t$  approaches 0. The set  $\widehat{W}$  is the set of endpoints of the paths that converge to a point in  $\mathbb{C}^N$  as  $t$  approaches 0.

One application of moving linear spaces is the homotopy membership test, first described in [21], which replaced the more expensive interpolation test of [20]. Given a point  $y \in \mathbb{C}^N$ , let  $\widehat{\mathcal{L}} : \mathbb{C}^N \rightarrow \mathbb{C}^\ell$  be a system of general linear polynomials such that  $y \in \mathcal{V}(\widehat{\mathcal{L}})$ . If  $\widehat{W}$  is the set of finite endpoints of the homotopy  $H$  defined in (1) starting at each point in  $W$ , then  $y \in V$  if and only if  $y \in \widehat{W}$ .

## 1.6 Geometric genus of a curve

In [3], a numerical algorithm is given for computing the geometric genus of an irreducible one-dimensional component  $R \subset \mathbb{C}^K$  of the solution set of a polynomial system. The geometric genus of  $R$  is the topological genus of the unique smooth compactification of the desingularization of  $R$ . Since the desingularizations of a curve and a generically one-to-one image of a curve are isomorphic, deflation of a component will not change its geometric genus. Therefore the component  $R$  may be assumed to have multiplicity one. The algorithm of [3], which is based on the Hurwitz theorem, starts with the restriction  $p : R \rightarrow \mathbb{C}$  of a linear projection  $A : \mathbb{C}^K \rightarrow \mathbb{C}$ .

In that article  $p$  is assumed proper, but this is *easily modified* as will be shown below. We also show how the algorithm may be applied to an irreducible curve  $R \subset \mathbb{C}^K$  arising as the closure of a constructible set  $R' \subset \mathbb{C}^K$ .

Let us explain the algorithm of [3].

We may regard  $A$  as the product projection  $\mathbb{C}^{K-1} \times \mathbb{C} \rightarrow \mathbb{C}$ . We let  $\overline{A}$  be the product projection of  $\mathbb{P}^{K-1} \times \mathbb{P}^1 \rightarrow \mathbb{P}^1$ . Taking the closure  $\overline{R}$  of  $R$  in  $\mathbb{P}^{K-1} \times \mathbb{P}^1$ , we have the proper map  $\overline{p} := \overline{A}_{\overline{R}}$ . Let  $s : \widehat{R} \rightarrow \overline{R}$  denote the desingularization of  $\overline{R}$  and  $\widehat{p} : \widehat{R} \rightarrow \mathbb{P}^1$

the map  $s \circ \bar{p}$ . Then, Hurwitz theorem tells us that

$$g = -2 \deg(\widehat{p}) + \rho$$

where

1.  $g$  is the genus of  $\widehat{R}$ , which we want to compute;
2.  $\deg(\widehat{p})$  is the degree of  $\widehat{p}$ , which equals the degree of  $p$ ; and
3.  $\rho$  is the ramification of  $\widehat{p}$ .

From the above we see that we need to compute  $\rho$ . Let  $\mathcal{R}$  denote the images under  $\widehat{p}$  of the branch points of  $\widehat{p}$ . For any  $y \in \mathcal{R}$ , let  $\Delta_y$  denote a contractible set with a continuous and piecewise differentiable boundary, e.g., a disk in a Euclidean patch  $\mathbb{C} \subset \mathbb{P}^1$  containing  $y$ , such that no points of  $\mathcal{R}$  other than  $y$  are in  $\Delta_y$ . Fix a point  $x_y$  on the boundary of  $\Delta_y$ . We have a monodromy transformation  $T_y : \widehat{p}^{-1}(x_y) \rightarrow \widehat{p}^{-1}(x_y)$  obtained by continuation around the boundary of  $\Delta_y$  of the paths starting at points of  $\widehat{p}^{-1}(x_y)$ . The ramification  $\rho$  is a sum of contributions  $\rho_y$  for the points  $y \in \mathcal{R}$ .

The number  $\rho_y$  equals  $\deg(\widehat{p})$  minus the number of orbits of the permutation group on  $\widehat{p}^{-1}(x_y)$  generated by  $T_y$ . There are two main observations of [3].

The first is that

- $\rho_y$  may be computed using the monodromy transformation  $T_y : \bar{p}^{-1}(x_y) \rightarrow \bar{p}^{-1}(x_y)$ . (We use the same symbol  $T_y$  because  $\bar{p}^{-1}(x_y)$  is naturally identified with  $\widehat{p}^{-1}(x_y)$  and under this identification, the monodromy transformations are the same.)

Though computing  $\mathcal{R}$  is involved, it is straightforward (see [3]) to compute a finite set on  $\bar{R}$  that maps to a finite set of  $\mathbb{P}^1$  containing  $\mathcal{R}$ . It suffices to work with this larger set instead of  $\mathcal{R}$  is a consequence of the second main observation:

- for any point  $y$  not in  $\mathcal{R}$ , the local monodromy contribution of  $\rho_y$  is zero.

Note also that we can work with  $p : R \rightarrow \mathbb{C}$  as long as we also do a calculation of  $\rho_\infty$  by going around a large enough circle on  $\mathbb{C}$ , so that any point of  $\mathcal{R}$  (except possibly for  $\infty$ ) is contained within the circle.

If  $p$  is not proper, we simply need to add to  $\mathcal{R}$  the points over which  $p$  is not proper, i.e., we add to  $\mathcal{R}$  the image under  $\bar{A}$  of the set  $\bar{R} \cap (\mathbb{P}^{K-1} \setminus \mathbb{C}^{K-1}) \times \mathbb{C}$ . As above, this may add a finite number of extra points without any harm to the final result.

**Extension to constructible sets** Finally, assume  $R' \subset \mathbb{C}^K$  is a constructible set whose closure is an irreducible curve  $R \subset \mathbb{C}^K$ . We fix a linear projection  $A : \mathbb{C}^K \rightarrow \mathbb{C}$  and set  $p$  equal to  $A$  restricted to  $R$ . Possibly making a linear change of coordinates, we regard  $A$  as the product projection  $\mathbb{C}^{K-1} \times \mathbb{C} \rightarrow \mathbb{C}$ . We let  $\bar{A}$  denote the product projection  $\mathbb{P}^{K-1} \times \mathbb{P}^1 \rightarrow \mathbb{P}^1$ . Let  $\bar{R}$  denote the closure of  $R$  in  $\mathbb{P}^{K-1} \times \mathbb{P}^1$ , we have the

proper map  $\bar{p} := \overline{A_{\bar{R}}}$ . We let  $\hat{p} : \hat{R} \rightarrow \mathbb{P}^1$  denote the composition of the desingularization map  $s : \hat{R} \rightarrow \bar{R}$  and  $\bar{p}$ .

Looking over the argument sketched above for the algorithm to compute  $g$  in the case when  $R' = R$ , we see that the algorithm for a constructible set  $R'$  to work, we need to compute

1. the degree of  $p$ ; and
2. a finite subset of  $\mathbb{P}^1$  containing  $\hat{p}(\mathcal{R})$ , where  $\mathcal{R}$  is the set of branch points of the  $\hat{p}$ .

The set  $\mathcal{R}$  is contained in the union of  $\infty \in \mathbb{P}^1$  and the images under  $p$  of

1. the singular points of  $R$ ;
2. the points  $R \setminus R'$ ;
3.  $A((\mathbb{P}^{K-1} \setminus \mathbb{C}^K) \cap \bar{R})$ ; and
4. all of the branch points of the algebraic map  $p : R \rightarrow \mathbb{C}$ .

## 2 Membership tests for projections

Let  $f : \mathbb{C}^N \rightarrow \mathbb{C}^n$  be a polynomial system and  $V \subset f^{-1}(0)$  be an irreducible algebraic set of dimension  $\ell$ . As developed in §1, we may assume without loss of generality that  $V$  is generically reduced and  $n = N - \ell$ .

Let  $\pi : \mathbb{C}^N \rightarrow \mathbb{C}^K$  be a linear map and  $y \in \pi(\mathbb{C}^N) \subseteq \mathbb{C}^K$ . We will first use a pseudo-witness set  $\{f, \pi, \mathcal{L}, \mathcal{W}\}$  for  $\overline{\pi(V)}$  to determine if  $y \in \overline{\pi(V)}$  and provide sufficient conditions for deciding if  $y \in \pi(V)$ . We will then use a witness set  $\{f, L, W\}$  for  $V$  to determine if  $y \in \pi(V)$ .

### 2.1 Basic membership test

Let  $\ell' = \dim \overline{\pi(V)}$  and  $\mathcal{L} = [\mathcal{L}_1 \cdots \mathcal{L}_{\ell'}]^T$  such that  $\mathcal{L}_1, \dots, \mathcal{L}_{\ell'}$  are general linear polynomials on  $\pi(\mathbb{C}^N)$  and  $\mathcal{L}_{\ell'+1}, \dots, \mathcal{L}_{\ell}$  are general linear polynomials on  $\mathbb{C}^N$ . For  $i = 1, \dots, \ell'$ , let  $\hat{\mathcal{L}}_i : \mathbb{C}^K \rightarrow \mathbb{C}$  be a general linear polynomial such that  $y \in \mathcal{V}(\hat{\mathcal{L}}_i)$  and define

$$\hat{\mathcal{L}}(x) = \begin{bmatrix} \hat{\mathcal{L}}_1(\pi(x)) \\ \vdots \\ \hat{\mathcal{L}}_{\ell'}(\pi(x)) \\ \mathcal{L}_{\ell'+1}(x) \\ \vdots \\ \mathcal{L}_{\ell}(x) \end{bmatrix}.$$

Consider the homotopy  $H$  defined by (1) which deforms  $\mathcal{L}$  to  $\widehat{\mathcal{L}}$ . For each  $w \in \mathcal{W}$ , let  $x_w(t)$  be the path defined by  $x_w(1) = w$  and  $H(x_w(t), t) \equiv 0$  for  $t \in (0, 1]$ . There are three possibilities for each path  $x_w(t)$  as  $t$  approaches 0, namely

1.  $x_w(t)$  converges to a point in  $\mathbb{C}^N$  yielding that  $\pi(x_w(t))$  converges to a point in  $\mathbb{C}^K$ ;
2.  $x_w(t)$  diverges but  $\pi(x_w(t))$  converges to a point in  $\mathbb{C}^K$ ; or
3.  $x_w(t)$  and  $\pi(x_w(t))$  both diverge.

Consider the related sets:

1.  $C_y = \{\lim_{t \rightarrow 0} \pi(x_w(t)) \mid w \in \mathcal{W} \text{ and } \lim_{t \rightarrow 0} x_w(t) \text{ converges}\}$ ; and
2.  $P_y = \{\lim_{t \rightarrow 0} \pi(x_w(t)) \mid w \in \mathcal{W} \text{ and } \lim_{t \rightarrow 0} \pi(x_w(t)) \text{ converges}\}$ .

Clearly,  $C_y \subset P_y \subset \overline{\pi(V)}$ . The following lemma yields a membership test for  $\overline{\pi(V)}$  using  $P_y$  and sufficient conditions for deciding if  $y \in \pi(V)$  using  $C_y$ .

**Lemma 1.** *With the setup described above, we have the following tests:*

1.  $y \in \overline{\pi(V)}$  if and only if  $y \in P_y$ ;
2. if  $y \in C_y$ , then  $y \in \pi(V)$ ; and
3. if  $C_y = P_y$  or  $\dim \overline{\pi(V)} = 1$ , then  $y \in C_y$  if and only if  $y \in \pi(V)$ .

*Proof.* Define  $\mathcal{L}_y(z) = \begin{bmatrix} \widehat{\mathcal{L}}_1(z) \\ \vdots \\ \widehat{\mathcal{L}}_{\ell'}(z) \end{bmatrix}$ . By genericity,  $\overline{\pi(V)} \cap \mathcal{V}(\mathcal{L}_y)$  consists of finitely many

points. It follows from [14] that  $P_y = \overline{\pi(V)} \cap \mathcal{V}(\mathcal{L}_y)$ . Since  $y \in \mathcal{V}(\mathcal{L}_y)$ , we know that  $y \in P_y$  if and only if  $y \in \overline{\pi(V)}$ .

If  $y \in C_y$ , then there exists  $w \in \mathcal{W}$  and  $\alpha \in V \subset \mathbb{C}^N$  such that  $\alpha = \lim_{t \rightarrow 0} x_w(t)$  and  $y = \pi(\alpha) = \lim_{t \rightarrow 0} \pi(x_w(t))$ . Since  $\alpha \in V$ , this implies  $y = \pi(\alpha) \in \pi(V)$ .

The only part remaining for the final statement is showing that  $y \in \pi(V)$  implies  $y \in C_y$ . If  $C_y = P_y$ , this follows from the first statement. If  $\dim \overline{\pi(V)} = 1$ , we know that  $y \in \pi(V)$  implies that  $V \cap \pi^{-1}(y)$  is pure-dimensional of dimension  $\dim V - 1$  since  $V$  is irreducible. Therefore, this case also follows from [14].  $\square$

**Remark 2.** *If  $w_1, w_2 \in \mathcal{W}$  such that  $\pi(w_1) = \pi(w_2)$ , then  $\pi(x_{w_1}(t)) = \pi(x_{w_2}(t))$  for all  $t \in (0, 1]$ . In particular, we only need to track at most  $\deg \pi(V) = |\pi(\mathcal{W})|$  paths in order to determine if  $y \in \overline{\pi(V)}$ .*

**Remark 3.** *We note that the membership test of this section immediately applies to a wide class of projections of quasialegbraic sets<sup>2</sup>. For example, consider the product*

<sup>2</sup>A quasialegbraic set is a set of the form  $A \setminus B$ , where  $A$  and  $B$  are algebraic subsets of  $\mathbb{P}^N$ .

projection  $\pi_{\mathbb{P}} : \mathbb{P}^{N-k} \times \mathbb{P}^K \rightarrow \mathbb{P}^K$ . Let  $X \subset \mathbb{P}^{N-k} \times \mathbb{P}^K$  be a quasialgebraic set. Let  $y \in \mathbb{P}^K$  be a point that we wish to check is in  $\overline{\pi_{\mathbb{P}}(X)}$ . Choose a generic Euclidean patch  $U \subset \mathbb{P}^{N-k} \times \mathbb{P}^K$ , i.e., choose generic hyperplanes  $H_v \subset \mathbb{P}^{N-k}$  and  $H_h \subset \mathbb{P}^K$  and let

$$U = (\mathbb{P}^{N-k} \setminus H_v) \times (\mathbb{P}^K \setminus H_h).$$

Then with probability one,  $y \in (\mathbb{P}^K \setminus H_h)$  and if  $y \in \overline{\pi_{\mathbb{P}}(X)} = \pi_{\mathbb{P}}(\overline{X})$ , i.e., if there is an  $x \in \overline{X}$  going to  $y$ , then  $x \in (\mathbb{P}^{N-k} \setminus H_v)$ .

## 2.2 Advanced membership test

We see from Lemma 1 that the one remaining case is deciding if  $y \in \pi(V)$  given that  $y \in P_y \subset \overline{\pi(V)}$ ,  $y \notin C_y$ , and  $\dim \overline{\pi(V)} > 1$ . The advanced membership test is based on the fact that  $y \in \pi(V)$  if and only if  $V \cap \pi^{-1}(y)$  is nonempty. That is, one simply computes the intersection of  $V$  with the linear space  $\pi^{-1}(y)$ . This can be accomplished starting with a witness set for  $V$  together with slice moving which we perform following a regenerative cascade approach [9]. Since we only need to decide if  $V \cap \pi^{-1}(y)$  is empty, the test simply cascades down through the dimensions under consideration and terminates when either a point in  $V \cap \pi^{-1}(y)$  is found or all of the possible dimensions are empty. As above, let  $\ell = \dim V$  and  $\ell' = \dim \overline{\pi(V)}$ . Then, since  $\ell - \ell'$  is the general fiber dimension, the possible fiber dimensions under consideration are  $\ell - 1, \ell - 2, \dots, \ell - \ell'$ . Thus, this test tracks at most  $\ell' \cdot \deg V$  paths. If we have already verified that  $y \notin C_y$  from §2.1, then we do not need to consider the general fiber dimension. In this case, this test tracks at most  $(\ell' - 1) \cdot \deg V$  additional paths.

Let  $\{f, L, W\}$  be a witness set for  $V$  where  $L = [L_1, \dots, L_{\ell}]^T$  and  $\widehat{\mathcal{L}}_1, \dots, \widehat{\mathcal{L}}_{\ell'}$  be general linear polynomials on  $\mathbb{C}^K$  such that  $y \in \mathcal{V}(\widehat{\mathcal{L}}_i)$ . For  $i = 0, \dots, \ell'$ , define

$$\mathcal{M}_i(x) = \begin{bmatrix} \widehat{\mathcal{L}}_1(\pi(x)) \\ \vdots \\ \widehat{\mathcal{L}}_i(\pi(x)) \\ L_{i+1}(x) \\ \vdots \\ L_{\ell}(x) \end{bmatrix}.$$

We have  $\mathcal{M}_0 = L$  and define  $S_0 = W$ . If  $0 \leq i < \ell'$  such that  $S_i$  is known, we compute  $S_{i+1}$  as follows. Let  $W_{i+1}$  be the finite endpoints of the modified homotopy  $H$  defined by (1) which deforms  $\mathcal{M}_i$  to  $\mathcal{M}_{i+1}$  with start points  $S_i$ . Let  $G_{i+1}$  be the subset of points of  $W_{i+1}$  which  $\pi$  maps to  $y$  and  $S_{i+1} = W_{i+1} \setminus G_{i+1}$ . In particular, it follows from [9, §2] that each point in  $S_{i+1}$  is a nonsingular root of  $\begin{bmatrix} f \\ \mathcal{M}_{i+1} \end{bmatrix}$  and  $G_{i+1}$  is a witness point superset for the pure  $(i+1)$ -codimensional component of  $V \cap \pi^{-1}(y)$ .



**Lemma 4.** *With the setup described above,  $y \in \pi(V)$  if and only if  $G_i \neq \emptyset$  for some  $i \in \{1, \dots, \ell'\}$ . Moreover, if  $y \notin C_y$ , where  $C_y$  is defined as in §2.1, then  $y \in \pi(V)$  if and only if  $G_i \neq \emptyset$  for some  $i \in \{1, \dots, \ell' - 1\}$ .*

*Proof.* This follows from the above discussion together with [9, Lemma 2.2 & Theorem 2.3] applied to this context.  $\square$

**Remark 5.** *By working with generic Euclidean patches as in Remark 3, the membership test of this section extends to a wide class of projections of quasiprojective sets. We will use the version for the restriction of the product projection  $\pi_{\mathbb{P}} : \mathbb{P}^{N-K} \times \mathbb{P}^K \rightarrow \mathbb{P}^K$  in the next section.*

### 3 Codimension one components of $\overline{\pi(V)} \setminus \pi(V)$

In order to compute more detailed invariants of  $\overline{\pi(V)}$ , it may be necessary to have a numerical irreducible decomposition of  $\overline{\pi(V)} \setminus \pi(V)$ , i.e., a numerical irreducible decomposition of  $\overline{\pi(V)} \setminus \pi(V)^\circ$ , where the set  $\mathcal{C}^\circ$  is the largest Zariski open set contained in the constructible algebraic set  $\mathcal{C}$ . The results of this article allow us to compute the decomposition of the codimension one components of  $\overline{\pi(V)} \setminus \pi(V)$ . As an illustration, we describe how to use this partial decomposition to compute a basic invariant of  $\overline{\pi(V)}$ .

#### 3.1 Decomposition of codimension one components

Assume we have the standard setup, i.e.,  $f : \mathbb{C}^N \rightarrow \mathbb{C}^n$  is a polynomial system and  $V \subset f^{-1}(0)$  is an irreducible  $\ell$ -dimensional component. By §1.2 and §1.3, we may assume without loss of generality that  $V$  is generically reduced and  $n = N - \ell$ , respectively. For simplicity, assume that  $\pi : \mathbb{C}^N \rightarrow \mathbb{C}^K$  is a linear projection onto the last  $K$  coordinates. Note that the projection  $\pi$  extends to the product projection  $\pi_{\mathbb{P}} : \mathbb{P}^{N-K} \times \mathbb{P}^K \rightarrow \mathbb{P}^K$ . Let  $V_{\mathbb{P}}$  denote the closure of  $V$  in  $\mathbb{P}^{N-K} \times \mathbb{P}^K$ .

Define  $\mathcal{P} = \mathbb{P}^{N-K} \times \mathbb{C}^K$  and  $\mathcal{E} = (\mathbb{P}^{N-K} \setminus \mathbb{C}^{N-K}) \times \mathbb{C}^K$ . Let  $V_{\mathcal{P}}$  denote the closure of  $V$  in  $\mathcal{P}$  and  $\pi_{\mathcal{P}}$  denote the restriction of  $\pi_{\mathbb{P}}$  to  $\mathcal{P}$ . We have the following:

1.  $\pi_{\mathbb{P}}(V_{\mathbb{P}})$  is the closure of  $\pi(V)$  in  $\mathbb{P}^K$ ; and
2.  $\overline{\pi(V)}$ , the closure of  $\pi(V)$  in  $\mathbb{C}^K$ , equals  $\pi_{\mathcal{P}}(V_{\mathcal{P}})$ .

We first consider the case when  $\dim \overline{\pi(V)} = 1$ . In this case, all fibers of  $\pi_{\mathcal{P}}$  restricted to  $V_{\mathcal{P}}$ , i.e.,  $\pi_{\mathcal{P}|V_{\mathcal{P}}}$ , are of pure dimension  $\dim \overline{\pi(V)} - 1$ . Since  $\dim(V_{\mathcal{P}} \cap \mathcal{E}) = \dim \overline{\pi(V)} - 1$ , we conclude that the irreducible components of the fibers of  $\pi_{\mathcal{P}|V_{\mathcal{P}}}$  over the finite set  $\overline{\pi(V)} \setminus \pi(V)$  are irreducible components of  $V_{\mathcal{P}} \cap \mathcal{E}$ . That is, we can first compute the numerical irreducible decomposition of  $V_{\mathcal{P}} \cap \mathcal{E}$  and then use the membership test of this article to determine the points of  $\overline{\pi(V)} \setminus \pi(V)$ .

Now assume that  $\dim \overline{\pi(V)} \geq 2$ . It is a consequence of a vanishing theorem of Picard-Kodaira type [18, Theorem 3.42] that, if  $\dim \overline{\pi(V)} \geq 2$ , then for a general hyperplane

$\mathcal{H}$  of  $\mathbb{C}^K$ ,  $H = \pi_{\mathcal{P}}^{-1}(\mathcal{H}) \cap V_{\mathcal{P}}$  is irreducible. Since a general linear space of codimension  $\dim \pi(V) + 1$  meets  $\overline{\pi(V)}$  in an irreducible curve and meets each codimension one component  $A$  of  $\overline{\pi(V)} \setminus \pi(V)$  in  $\deg A$  points, we have constructed a pseudo-witness point set for the union of codimension one components of  $\overline{\pi(V)} \setminus \pi(V)$ . This shows in particular that for each codimension one irreducible component  $A$  of  $\overline{\pi(V)} \setminus \pi(V)$ , there is a  $\dim V - 1$  component of  $V_{\mathcal{P}} \cap \mathcal{E}$  surjecting onto  $A$ .

We may compute a numerical irreducible decomposition of the  $\dim \overline{V} - 1$  components  $\mathcal{A}$  of  $V_{\mathcal{P}} \cap \mathcal{E}$  and then use the membership test for the images of the witness sets of these components to check which components have images in  $\overline{\pi(V)} \setminus \pi(V)$ . For the irreducible components  $\mathcal{A}$  with an image  $A$  in  $\overline{\pi(V)} \setminus \pi(V)$ , the results of [7] using the map  $\pi_{\mathcal{A}} : \mathcal{A} \rightarrow A$  yield a pseudo-witness set of  $A$ . Finally, we need to compute the dimensions of the fibers over the images of one point from the witness set of  $\mathcal{A}$ : those  $A$  of codimension one are precisely the ones where the fiber dimension is  $\dim V - \dim \pi(V)$ .

### 3.2 Application to the geometric genus of a curve section

As an application, we describe how we may use the pseudo-witness set for the codimension one boundary components to compute the geometric genus  $g$  of a general curve section of  $\overline{\pi(V)} \subset \mathbb{C}^K$ . The number  $g$  is, by definition, equal to the genus of the desingularization  $s : \widehat{R} \rightarrow R$  of the intersection  $R \subset \mathbb{C}^K$  of  $\overline{\pi(V)}$  and a general affine linear space  $L$  of  $\mathbb{C}^K$  of dimension  $K + 1 - \dim \overline{\pi(V)}$ . Topologically,  $g$  is the usual genus of the unique smooth compactification of  $\widehat{R}$ . Equivalently, it is the genus of the desingularization of the closure in  $\mathbb{P}^K$  of the intersection of  $\overline{\pi(V)}$  and a general affine linear space of  $\mathbb{C}^K$  of dimension  $K + 1 - \dim \overline{\pi(V)}$ .

Given a general affine linear space  $L$  of  $\mathbb{C}^K$  of dimension  $K + 1 - \dim \overline{\pi(V)}$ , we take  $R = \overline{\pi(V)} \cap L$  and  $R' = \pi(V) \cap L$ . Using [18, Theorem 3.42], we can again reduce down to the case that  $\dim \pi(V) = 1$ . We take the map  $p : R \rightarrow \mathbb{C}$  to be the restriction to  $R$  to any linear projection from  $\mathbb{C}^K$  to  $\mathbb{C}$ . By taking the intersection of  $V$  with a general affine linear space of codimension  $\ell - 1$ , we may, by renaming if necessary, assume that  $V$  is one-dimensional. Let  $q : V \rightarrow R$  be the map obtained by composing the restriction of  $\pi$  to  $V$  with the map  $p$ .

For  $Q$ , we take the union of the following sets:

1. the image under  $q$  of the branchpoints of  $q$ ;
2. the image under  $q$  of the singular points of  $V$ ;
3. the image under  $p \circ \pi_{\mathcal{P}}$  of the set  $V_{\mathcal{P}} \cap \mathcal{E}$  using the notation from §3.1; and
4. the image under  $p$  of the points in  $\overline{\overline{\pi(V)} \setminus \pi(V)}$ .

The first three items require only standard computations. The last item follows from the computation of the finite set  $\overline{\pi(V)} \setminus \pi(V)$ , which was computed in §3.1. Note the

third item is a finite set of points containing the points over which  $q$  is not proper and therefore also the points over which  $p$  is not proper.

To see that this set  $Q$  suffices, note that all the singular points of  $\overline{\pi(V)}$  and branchpoints of  $p$  are either over  $\overline{\pi(V)} \setminus \pi(V)$  or in the image of the branchpoints of  $p \circ \pi_{\mathcal{P}}$  and the singular set of  $V_{\mathcal{P}}$ .

The last item needed is the ability to track paths. We note that the paths on  $\overline{\pi(V)}$  which are contained in  $\pi(V)^{\circ}$  may be tracked using the pseudo-witness set of  $\overline{\pi(V)}$ .

## 4 Examples

We conclude by demonstrating the membership tests and codimension one decomposition on illustrative examples and then report on a more advanced example. The linear slice moving computations reported here were performed using Bertini v1.3.1 [2].

### 4.1 A parameterized circle

Consider the rational parameterization  $(x(s), y(s)) = \left( \frac{1-s^2}{1+s^2}, \frac{2s}{1+s^2} \right)$  of an open dense subset the unit circle. Clearing denominators, this parameterization yields the system

$$f(s, x, y) = \begin{bmatrix} x(1+s^2) - (1-s^2) \\ y(1+s^2) - 2s \end{bmatrix}$$

with the accompanying projection  $\pi(s, x, y) = (x, y)$  defined by the matrix  $B = [0 \ I_2]$  where  $I_2$  is the  $2 \times 2$  identity matrix. It is easy to verify that  $V = f^{-1}(0)$  is a irreducible curve of degree 3 that is generically reduced with respect to  $f$ . We will first use a witness set for  $V$  to construct a pseudo-witness set for  $\overline{\pi(V)}$  and then determine if  $z_j \in \pi(V)$  and  $z_j \in \overline{\pi(V)}$  where

$$z_1 = (0, 1), \quad z_2 = (-1, 0), \quad z_3 = (\sqrt{2}, i), \quad \text{and} \quad z_4 = (1 + i, 1/3 - i/2)$$

with  $i = \sqrt{-1}$ . Finally, we will compute  $\overline{\pi(V)} \setminus \pi(V)$ .

**Pseudo-witness set construction:** Let  $\{f, L, W\}$  be a witness set for  $V$  where  $L : \mathbb{C}^3 \rightarrow \mathbb{C}$  is a general linear polynomial and  $|W| = 3$ . Since, for any  $w \in W$ ,

$$\begin{bmatrix} Jf(w) \\ B \end{bmatrix}$$

is full rank, where  $Jf(w)$  is the Jacobian matrix of  $f$  evaluated at  $w$ , Lemma 3 of [7] yields that  $\dim \pi(V) = 1$ . Let  $\mathcal{L}(s, x, y) = \alpha x + \beta y - 1$  where  $\alpha, \beta \in \mathbb{C}$  are random, which is a linear polynomial in the image of  $\pi$ . Consider the three paths defined by modifying the homotopy  $H$  from (1) to move from  $L$  to  $\mathcal{L}$  starting at the three points in  $W$ . Two paths converge with their endpoints mapping to distinct points under  $\pi$ . This implies that the degree of  $\overline{\pi(V)}$  is 2. If  $\mathcal{W}$  is the set consisting of these two endpoints, then  $\{f, \pi, \mathcal{L}, \mathcal{W}\}$  is a pseudo-witness set for  $\overline{\pi(V)}$ .

$j$	$ C_{z_j} $	$z_j \in C_{z_j}?$	$ P_{z_j} $	$z_j \in P_{z_j}?$	Result from Lemma 1
1	2	Yes	2	Yes	$z_1 \in \pi(V)$
2	1	No	2	Yes	$z_2 \in \overline{\pi(V)} \setminus \pi(V)$
3	2	Yes	2	Yes	$z_3 \in \pi(V)$
4	2	No	2	No	$z_4 \notin \pi(V)$

Table 1: Summary of basic membership test for unit circle

**Basic membership test:** For each  $j = 1, \dots, 4$ , let  $z_j = (z_j^x, z_j^y)$  and consider the linear polynomial  $\widehat{\mathcal{L}}_j(s, x, y) = \alpha(x - z_j^x) + \beta(y - z_j^y)$ . The basic membership test described in §2.1 uses a modification of the homotopy  $H$  from (1) to move from  $\mathcal{L}$  to  $\widehat{\mathcal{L}}_j$  starting with the two points in  $\mathcal{W}$ . Since  $\dim \overline{\pi(V)} = 1$ , Lemma 1 provides membership tests for both  $\pi(V)$  and  $\overline{\pi(V)}$ . Table 1 summarizes the results. Here, the sets  $C_{z_j}$  and  $P_{z_j}$  are the sets as in §2.1 arising from this basic membership test.

**Codimension one components:** The codimension one components of  $\overline{\pi(V)} \setminus \pi(V)$  correspond to the points in  $\pi_{\mathcal{P}}(V_{\mathcal{P}} \cap \mathcal{E})$  (as defined in §3.1). By working on a random patch in  $\mathbb{P}^1$ , this reduces to tracking paths in  $\mathbb{C}^4$ . We homogenize  $f$  and  $L$  with respect to  $s$ , namely

$$F(s_0, s_1, x, y) = s_0^2 \cdot f\left(\frac{s_1}{s_0}, x, y\right) \quad \text{and} \quad M(s_0, s_1, x, y) = s_0 \cdot L\left(\frac{s_1}{s_0}, x, y\right),$$

and fix an affine patch in  $\mathbb{P}^1$  defined by the equation  $P(s_0, s_1, x, y) = p_0 s_0 + p_1 s_1 - 1$  where  $p_0, p_1 \in \mathbb{C}$  are random. Let  $\widehat{M}$  be a general linear form and  $M_0(s_0, s_1, x, y) = s_0$ . Starting with the points

$$S = \left\{ \left( \frac{1}{p_1 + p_2 s}, \frac{s}{p_1 + p_2 s}, x, y \right) \mid (s, x, y) \in W \right\},$$

we first compute the finite endpoints of the homotopy  $H$  from (1) modified to deform from  $[M, P]^T$  to  $[\widehat{M}, P]^T$  and then use those as start points as we deform from  $[\widehat{M}, P]^T$  to  $[M_0, P]^T$ . The resulting finite endpoints correspond to the points in  $V_{\mathcal{P}} \cap \mathcal{E}$ . In particular, the first had three finite endpoints while only one of the three paths converged for the second. This endpoint corresponds to the point  $(0, 1, -1, 0) \in \mathbb{P}^1 \times \mathbb{C}^2$ . Since this point projects to  $(-1, 0)$  under  $\pi_{\mathcal{P}}$  (as defined in §3.1), we know that  $\overline{\pi(V)} \setminus \pi(V) = \{(-1, 0)\}$ .

## 4.2 A two-dimensional constructible set

Consider the example from §1.4, namely the image of  $V = f^{-1}(0)$  under the projection  $\pi(s, x, y) = (x, y)$  where  $f(s, x, y) = x - sy$ . The projection  $\pi$  is defined by the matrix  $B = [0 \ I_2]$  where  $I_2$  is the  $2 \times 2$  identity matrix. Clearly,  $V$  is an irreducible surface of degree 2 that is generically reduced with respect to  $f$ . After constructing a pseudo-witness set for  $\overline{\pi(V)}$ , we will use the membership tests to determine if  $p_j \in \pi(V)$  where

$$p_1 = (1, 1), \quad p_2 = (0, 0), \quad \text{and} \quad p_3 = (1, 0),$$

$j$	$ C_{p_j} $	$p_j \in C_{p_j}?$	$ P_{p_j} $	$p_j \in P_{p_j}?$	Result from Lemma 1
1	1	Yes	1	Yes	$p_1 \in \pi(V)$
2	1	Yes	1	Yes	$p_2 \in \pi(V)$
3	0	No	1	Yes	$p_3 \in \overline{\pi(V)}$ , inconclusive on $\pi(V)$

Table 2: Summary of basic membership test

and then compute a decomposition of the codimension one components of  $\overline{\pi(V)} \setminus \pi(V)$ .

**Pseudo-witness set construction:** Let  $\{f, L, W\}$  be a witness set for  $V$  where  $L : \mathbb{C}^3 \rightarrow \mathbb{C}^2$  is a system of general linear polynomials and  $|W| = 2$ . Since, for any  $w \in W$ ,

$$\begin{bmatrix} \nabla f(w)^T \\ B \end{bmatrix}$$

is full rank, where  $\nabla f(w)$  is the gradient of  $f$  evaluated at  $w$ , Lemma 3 of [7] yields that  $\dim \overline{\pi(V)} = 2$ . Therefore,  $\overline{\pi(V)} = \mathbb{C}^2$  and  $\deg \overline{\pi(V)} = 1$ . For random  $\alpha, \beta \in \mathbb{C}$ , let

$$\mathcal{L}(s, x, y) = \begin{bmatrix} x - \alpha \\ y - \beta \end{bmatrix}.$$

A pseudo-witness set for  $\overline{\pi(V)}$  is the quadruple  $\{f, \pi, \mathcal{L}, \mathcal{W}\}$  where  $\mathcal{W} = \{(\alpha/\beta, \alpha, \beta)\}$ .

**Basic membership test:** Even though  $\overline{\pi(V)} = \mathbb{C}^2$  and hence  $p_j \in \overline{\pi(V)}$ , we can still use the basic membership test of §2.1 to determine which points to further investigate using the advanced membership test of §2.2. For each  $j = 1, 2, 3$ , we used the system of linear polynomials  $\widehat{\mathcal{L}}_j(s, x, y) = (x, y) - p_j$ . Table 2 summarizes the results. Here, the sets  $C_{p_j}$  and  $P_{p_j}$  are the sets as in §2.1 arising from this basic membership test.

**Advanced membership test:** Since the basic membership test was inconclusive for deciding if  $p_3 = (1, 0) \in \pi(V)$ , we now apply the advanced membership test of §2.2. Let  $L = [L_1, L_2]^T$  where  $L$  is the linear system in the witness set  $\{f, L, W\}$  for  $V$ . For  $i = 1, 2$ , let  $\widehat{\mathcal{L}}_i(s, x, y) = r_{i1}(x - 1) + r_{i2}y$  for random  $r_{ij} \in \mathbb{C}$  and consider

$$\mathcal{M}_0 = \begin{bmatrix} L_1 \\ L_2 \end{bmatrix}, \quad \mathcal{M}_1 = \begin{bmatrix} \widehat{\mathcal{L}}_1 \\ L_2 \end{bmatrix}, \quad \text{and} \quad \mathcal{M}_2 = \begin{bmatrix} \widehat{\mathcal{L}}_1 \\ \widehat{\mathcal{L}}_2 \end{bmatrix}.$$

Starting with  $S_0 = W$ , tracking the paths for the modified homotopy  $H$  from (1) that deforms from  $\mathcal{M}_0$  to  $\mathcal{M}_1$  produces two points, neither of which projects to  $p_3$ . Since we have already performed the basic membership test and found that  $p_3 \notin C_{p_3}$ , Lemma 4 provides that  $p_3 \notin \pi(V)$ . If the basic test was not already performed, one would need to track the two paths arising from moving  $\mathcal{M}_1$  to  $\mathcal{M}_2$ . Since both of these paths diverge, the same conclusion is reached.

Since the endpoint of the path for  $p_2 = (0, 0)$  was singular when performing the basic membership test, it is instructive to perform the advanced membership test on this point as well. In this case, we take  $\widehat{\mathcal{L}}_i(s, x, y) = r_{i1}x + r_{i2}y$ . The deformation from  $\mathcal{M}_0$  to

$\mathcal{M}_1$  also produces two points, one of which does project to  $p_2$ . Therefore, we know that  $V \cap \pi^{-1}(p_2)$  contains a line. Tracking from the other point as  $\mathcal{M}_1$  moves to  $\mathcal{M}_2$  produces another point on this line. Therefore,  $V \cap \pi^{-1}(p_2)$  is a line, namely  $\{(s, 0, 0) \mid s \in \mathbb{C}\}$ .

**Codimension one components:** We now turn to computing the curves in  $\mathbb{C}^2$  contained in  $\overline{\pi(V)} \setminus \pi(V)$  which correspond to the curves in  $\pi_{\mathcal{P}}(V_{\mathcal{P}} \cap \mathcal{E})$  (as defined in §3.1). As in §4.1, we perform this computation on a random patch in  $\mathbb{P}^1$  which reduces to tracking paths in  $\mathbb{C}^4$ . We homogenize  $f$  and  $L$  with respect to  $s$ , namely

$$F(s_0, s_1, x, y) = s_0 \cdot f\left(\frac{s_1}{s_0}, x, y\right) = s_0x - s_1y \quad \text{and} \quad M(s_0, s_1, x, y) = s_0 \cdot L\left(\frac{s_1}{s_0}, x, y\right),$$

and fix an affine patch in  $\mathbb{P}^1$  defined by the equation  $P(s_0, s_1, x, y) = p_0s_0 + p_1s_1 - 1$  where  $p_0, p_1 \in \mathbb{C}$  are random. Let  $\widehat{M} = [\widehat{M}_1, \widehat{M}_2]^T$  be a system of two general linear forms and  $M_0(s_0, s_1, x, y) = [s_0, \widehat{M}_2(s_0, s_1, x, y)]^T$ . Starting with the points

$$S = \left\{ \left( \frac{1}{p_1 + p_2s}, \frac{s}{p_1 + p_2s}, x, y \right) \mid (s, x, y) \in W \right\},$$

we first compute the finite endpoints of the homotopy  $H$  from (1) modified to deform from  $[M, P]^T$  to  $[\widehat{M}, P]^T$  and then use those as start points as we deform from  $[\widehat{M}, P]^T$  to  $[M_0, P]^T$ . We can use the resulting finite endpoints to produce a witness set for the curves in  $V_{\mathcal{P}} \cap \mathcal{E}$ . In particular, the first had two finite endpoints while only one of the two paths converged for the second. This endpoint corresponds to the point  $(0, 1, x^*, 0) \in \mathbb{P}^1 \times \mathbb{C}^2$  where  $x^* \in \mathbb{C}$ . This point projects to  $(x^*, 0)$  under  $\pi_{\mathcal{P}}$  (as defined in §3.1) which forms a witness point set for the line in  $\overline{\pi(V)} \setminus \pi(V)$ , namely  $\mathcal{V}(y)$ .

### 4.3 Lüroth hypersurface

Classically, a plane quartic is called a Lüroth quartic if it contains the ten vertices of a complete pentilateral [13]. The closure of the set of classical Lüroth quartics is a hypersurface  $\mathcal{H}$  in the space of plane quartics, called the Lüroth hypersurface, that was showed by Morley in 1919 to have degree 54 [15]. The degree 54 polynomial equation defining this hypersurface is called the Lüroth invariant and, according to Ottaviani [16], it is still unknown. Without the defining equation, deciding if a given quartic lies on the Lüroth hypersurface requires another approach. The approach in [16] provides a partial test which builds on classical results of White and Miller [24]. We will use a pseudo-witness set for  $\mathcal{H}$  and the membership test of §2.1 to provide a complete test for deciding if a given plane quartic lies on  $\mathcal{H}$  by tracking at most 54 homotopy paths.

**Pseudo-witness set construction:** We construct a pseudo-witness set for  $\mathcal{H}$  by first identifying the space of plane quartics with  $\mathbb{P}^{14}$  so that  $\mathcal{H} \subset \mathbb{P}^{14}$ . The set  $\mathcal{H}$  is the closure of the set of plane quartics  $Q$  for which there exists nonzero linear polynomials

$j$	$ C_{Q_j} $	$Q_j \in C_{Q_j}?$	$ P_{Q_j} $	$Q_j \in P_{Q_j}?$	Result from Lemma 1
1	54	No	54	No	$Q_1 \notin \mathcal{H}$
2	54	No	54	No	$Q_2 \notin \mathcal{H}$
3	54	No	54	No	$Q_3 \notin \mathcal{H}$
4	54	No	54	No	$Q_4 \notin \mathcal{H}$
5	54	Yes	54	Yes	$Q_5 \in \mathcal{H}$
6	38	No	39	Yes	$Q_6 \in \mathcal{H}$

Table 3: Summary of membership in the hypersurface of Lüroth quartics

$\ell_j$  for  $j = 1, \dots, 5$  such that

$$Q = \mathcal{V} \left( \sum_{j=1}^5 \prod_{\substack{k=1 \\ k \neq j}}^5 \ell_k \right).$$

This parameterization allows us to use Lemma 3 of [7] to confirm that  $\mathcal{H}$  is a hypersurface and compute a pseudo-witness set for  $\mathcal{H}$  using Bertini [2]. From this pseudo-witness set, we are able to verify Morley's result that the degree of  $\mathcal{H}$  is 54 and, following Remark 2, we chose 54 points from the pseudo-witness point set that correspond to distinct quartics to be used as the starting points for our basic membership test.

**Basic membership test:** We applied the basic membership test of §2.1 to the quartics  $Q_j = \mathcal{V}(q_j)$  defined by the following polynomials:

- $q_1 = (x^2 + y^2 + z^2)^2$ ;
- (Edge quartic [6, 17])  $q_2 = 25(x^4 + y^4 + z^4) - 34(x^2y^2 + x^2z^2 + y^2z^2)$ ;
- (Klein quartic [16, §5])  $q_3 = x^3y + y^3z + z^3x$ ;
- (Vinnikov curve [17, Ex. 4.1])  $q_4 = 2x^4 + y^4 + z^4 - 3x^2y^2 - 3x^2z^2 + y^2z^2$ ;
- ([16, §5])  $q_5 = xyz(x + y + z) + (x + 2y + 3z)(xyz + (xy + xz + yz)(x + y + z))$ ; and
- $q_6 = x^3y + x^2z^2 + xz^3$ .

Table 3 summarizes the results of this test. Here, the sets  $C_{Q_j}$  and  $P_{Q_j}$  are the sets as in §2.1 arising from this basic membership test. We note that since  $\mathcal{H} \subset \mathbb{P}^{14}$ , compactness yields that  $P_{Q_j}$ , as a list, must consist of 54 points. However, in the  $j = 6$  case, 16 of these points coincided with  $Q_6$ .

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